

COMPUTATIONAL COMPLEXITY THEORY (COMS 30126) GUIDE TO EXAMINABLE MATERIAL

★ Anything not explicitly listed as omitted should be regarded as examinable. (This includes unlisted topics). If in doubt, please enquire!

★ For many theorems, *proofs* are omitted, but *statement* of the theorem is still required i.e. you need to be able to accurately state the theorem and understand what it says!

★ All exercise sheets are examinable material with the following questions OMITTED: sheet 1: 1,4,7; sheet 2: 5; sheet 3: 3; sheet 4: 3; sheet 5: no omissions; sheet 6: 3,4.

COURSE OUTLINE WITH OMITTED ITEMS IDENTIFIED:

TMs: basic structure and operation (transition table). How to specify a configuration.

NDTMs: definition, computational paths, branching tree of all paths, notion of accepting/rejecting an input (asymmetrical definition).

Languages $\mathcal{L} = \{0, 1\}^*$. TM or NDTM recognising/deciding a language.

Theorem: set of all TMs can be enumerated. OMIT proof.

Theorem: halting problem is undecidable. OMIT proof.

Big/small O notation. Statements of principal properties comparing log/poly/exponential growths.

Theorem (page 12): NDTMs can be simulated on TM with exponential time overhead. OMIT proof.

The class P.

PATH problem. Adjacency matrix of a directed graph. Argument that $\text{PATH} \in \text{P}$.

Verifiers. Poly-time verifiers.

The class NP (defined by verifiers). co-NP.

Theorem: $A \in \text{NP}$ iff A is decided by a poly-time NDTM. OMIT proof.

Notion of poly-time reduction. NP completeness.

Satisfiability problem. Cook-Levin theorem. OMIT proof.

Theorem: 3-SAT is NPC. OMIT proof.

Oracle TMs (page 18-20): definition of oracle TM. Theorems A and B. OMIT proofs.

Theorem: 2-SAT \in P. OMIT proof.

Space complexity. PSPACE.

Savitch theorem. OMIT proof.

$\text{NPSPACE} = \text{PSPACE}$ and proof.

PSPACE completeness, definition.

Definitions of: space bounded TMs, log space transducer and its configurations. L, NL, NL completeness.

Theorem: PATH is NL complete. OMIT proof.

Randomised complexity: OMIT examples of randomised algorithms (on page 27-29).

Definitions of: PTMs, accept/reject probability for an input, BPP.

Amplification lemma statement. OMIT proof.

Definitions of: RP, co-RP, ZPP, PP, MAJSAT.

Theorem: $NP \subseteq PP$ and proof.

Circuits: definition of Boolean circuit.

Proposition (page 35): universality of AND, OR, NOT gates, and proof.

Theorem: “most Boolean functions need exp big circuits”. OMIT proof.

Definitions of: circuit family, decision of \mathcal{L} by a circuit family, size and depth complexity.

Example of parity function.

Definition of uniform circuit family.

Proposition: “undecidable languages can have poly sized circuits”. and proof.

Theorem C: $\mathcal{L} \in P$ iff \mathcal{L} has uniform circuit family. OMIT proof.

Notion of P completeness. CIRCUIT-VALUE problem.

Theorem: CIRCUIT-VALUE is P complete, and proof.

Theorem: any $\mathcal{L} \in BPP$ has poly sized circuits. OMIT proof.

Definition of parallel complexity classes NC^i and NC.

Theorem: $NC^1 \subseteq L \subseteq NL \subseteq NC^2$. OMIT proofs.

Definition of interactive proof system and class IP.

$BPP \subseteq IP$, $NP \subseteq IP$.

Remarks on probabilities in the definition: OMIT these remarks.

Theorem: $dIP = NP$, and proof.

Graph non-isomorphism problem and its interactive proof system.

Program checking application: OMIT this section.

Theorem: $IP = PSPACE$. OMIT proof.

Theorem: $\#SAT$ is in IP. OMIT proof.